

СТАС КЕЛЬВИЧ, POSTGRES PRO

РАСПРЕДЕЛЕННЫЕ ТРАНЗАКЦИИ И ПУТЕШЕСТВИЯ ВО ВРЕМЕНИ

We believe it is better to have application programmers deal with performance problems due to overuse of transactions as bottlenecks arise, rather than always coding around the lack of transactions.

оригинальная статья о Google Spanner

ПОГОВОРИМ:

- ▶ Сериализуемость
- ▶ Однопоточная сериализуемость
- ▶ Двухфазные локи, 2PL
- ▶ Snapshot isolation
- ▶ Распределенный snapshot isolation
- ▶ Clock-SI

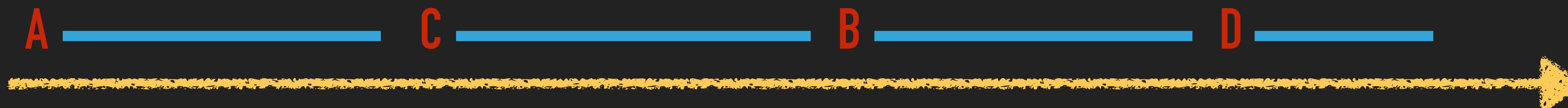
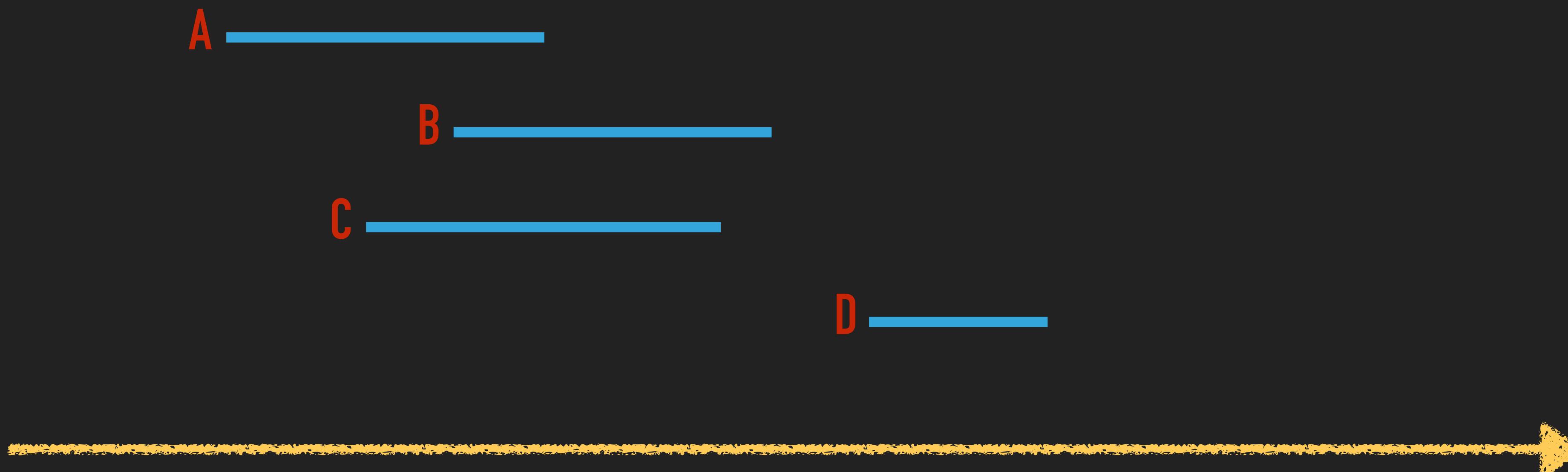
ATOMICITY
CONSISTENCY
ISOLATION ->
DURABILITY



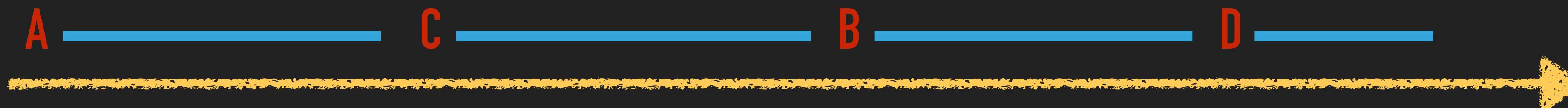
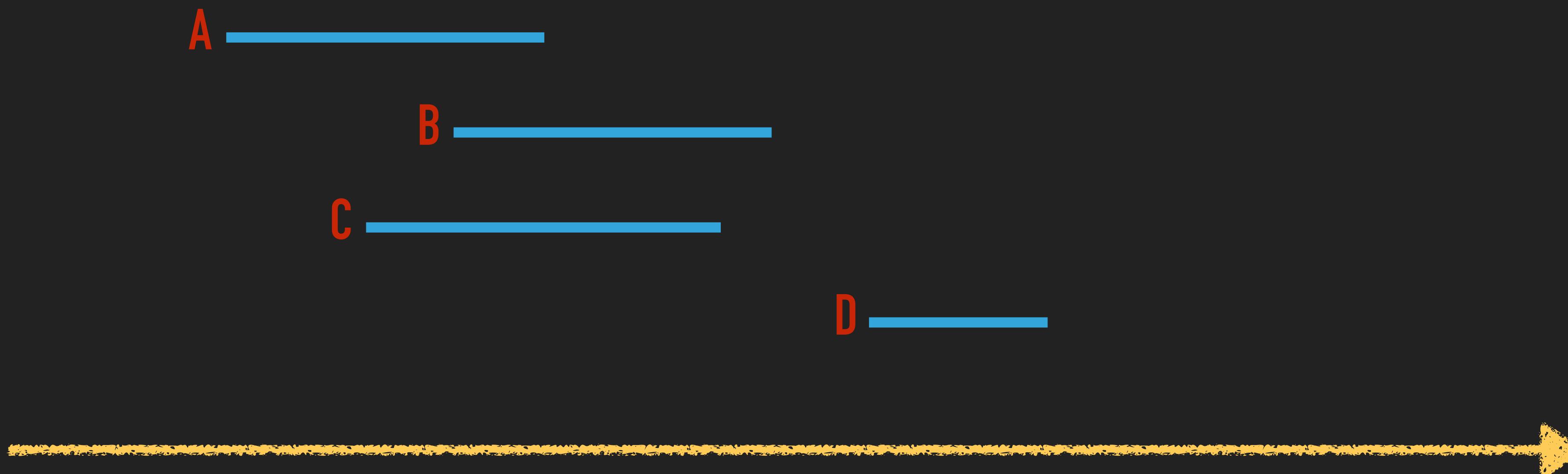
Вас много, а я одна.

База данных о клиентских подключениях

СЕРИАЛИЗУЕМОСТЬ



СЕРИАЛИЗУЕМОСТЬ



50 ОТТЕНКОВ СЕРИАЛИЗАЦИИ

VSR OCSR COCSR 1SR
SS2PL FSR CMFSR SSR
CMVSR CSR

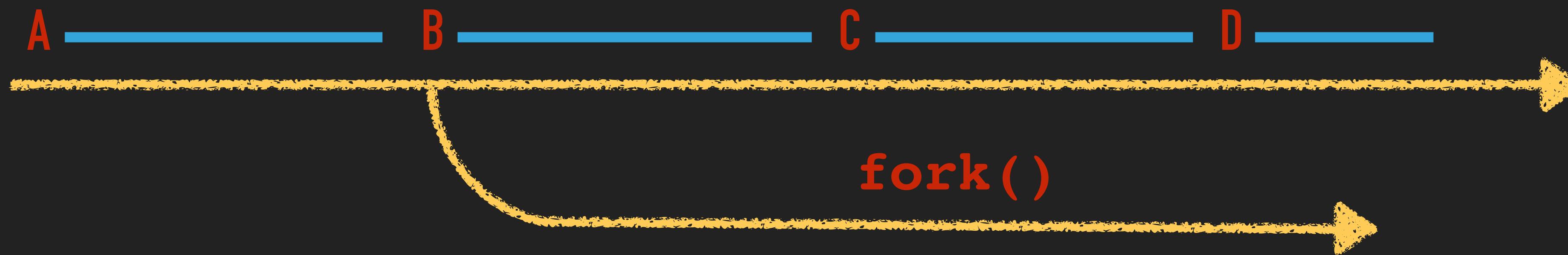
СЕРИАЛИЗУЕМОСТЬ В ОДНОМ ПОТОКЕ

REDIS, TARANTOOL



СЕРИАЛИЗУЕМОСТЬ В ОДНОМ ПОТОКЕ

REDIS, TARANTOOL



УРОВНИ ИЗОЛЯЦИИ И 2PL

READ UNCOMMITTED

Без read-blokirovki

READ COMMITTED

Короткая read-blokirovka

REPEATABLE READ

Длинная read-blokirovka

SERIALIZABLE

Длинная read-blokirovka + !фантомы

УРОВНИ ИЗОЛЯЦИИ И 2PL

A Critique of ANSI SQL Isolation Levels

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Abstract: ANSI SQL-92 [MS, ANSI] defines Isolation Levels in terms of *phenomena*: Dirty Reads, Non-Repeatable Reads, and Phantoms. This paper shows that these phenomena and the ANSI SQL definitions fail to properly characterize several popular isolation levels, including the standard locking implementations of the levels covered. Ambiguity in the statement of the phenomena is investigated and a more formal statement is arrived at; in addition new phenomena that better characterize isolation types are introduced. Finally, an important multiversion isolation type, called Snapshot Isolation, is defined.

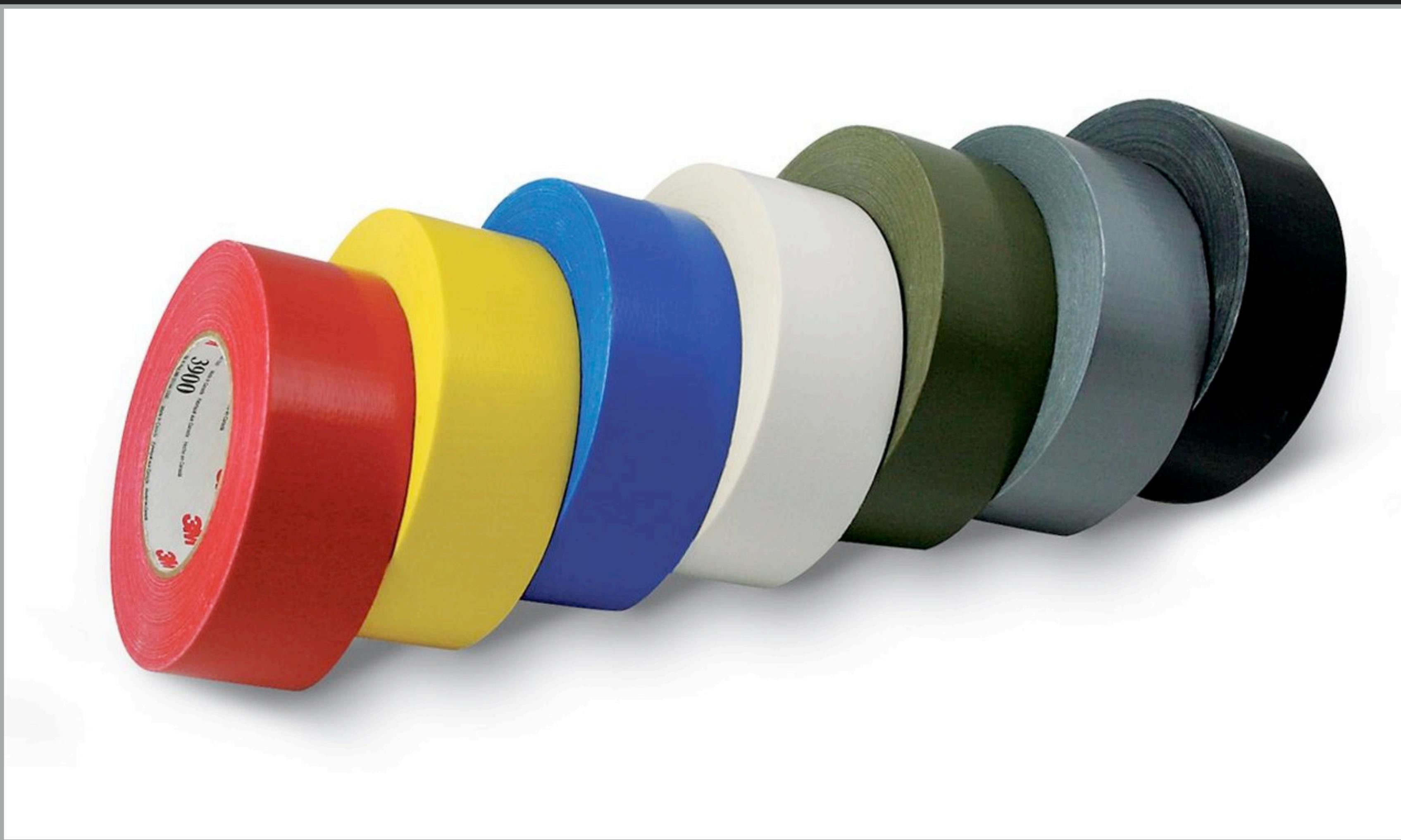
1 . Introduction

Running concurrent transactions at different isolation levels allows application designers to trade off concurrency and throughput for correctness. Lower isolation levels increase

The ANSI isolation levels are related to the behavior of lock schedulers. Some lock schedulers allow transactions to vary the scope and duration of their lock requests, thus departing from pure two-phase locking. This idea was introduced by [GLPT], which defined *Degrees of Consistency* in three ways: locking, data-flow graphs, and anomalies. Defining isolation levels by phenomena (anomalies) was intended to allow non-lock-based implementations of the SQL standard.

This paper shows a number of weaknesses in the anomaly approach to defining isolation levels. The three ANSI phenomena are ambiguous, and even in their loosest interpretations do not exclude some anomalous behavior that may arise in execution histories. This leads to some counter-intuitive results. In particular, lock-based isolation levels have different characteristics than their ANSI equivalents.

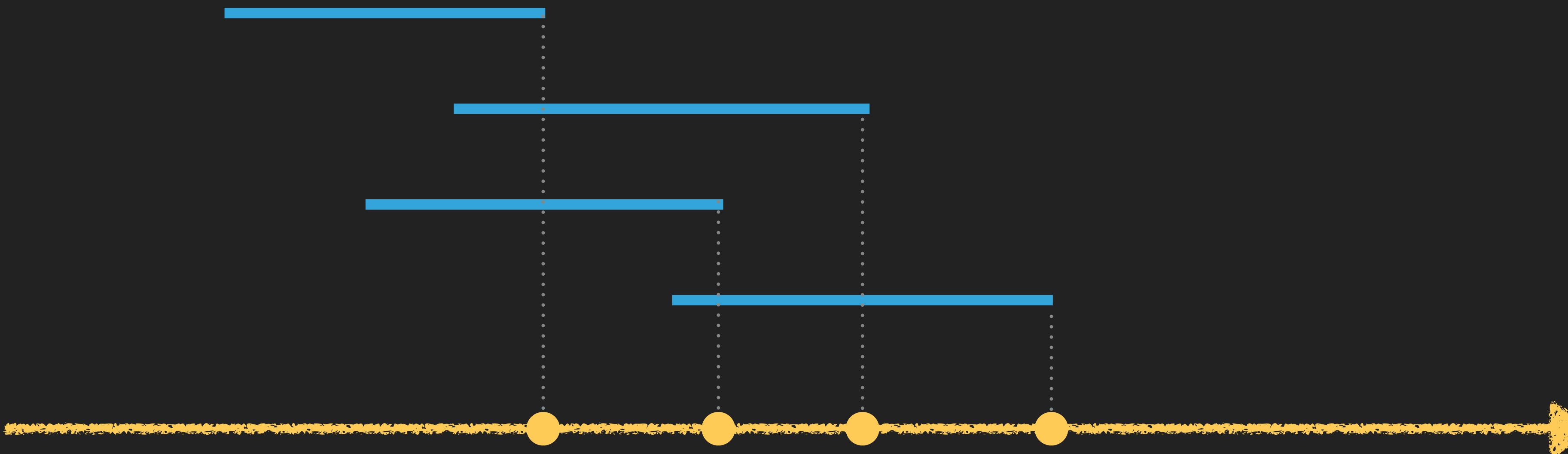
УРОВНИ ИЗОЛЯЦИИ И 2PL



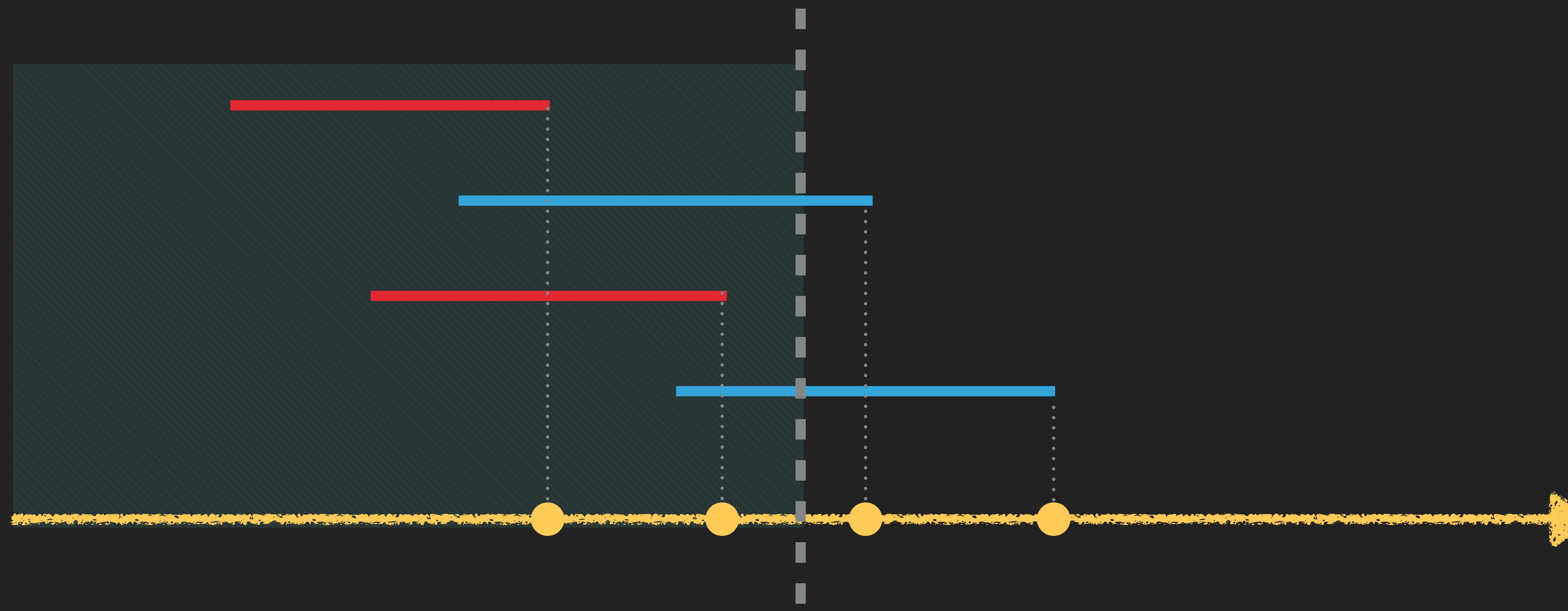
SNAPSHOT ISOLATION

- ▶ Не изменяем данные *in-place*, а создаем версии
- ▶ Читатели не блокируют писателей, писатели не блокируют читателей

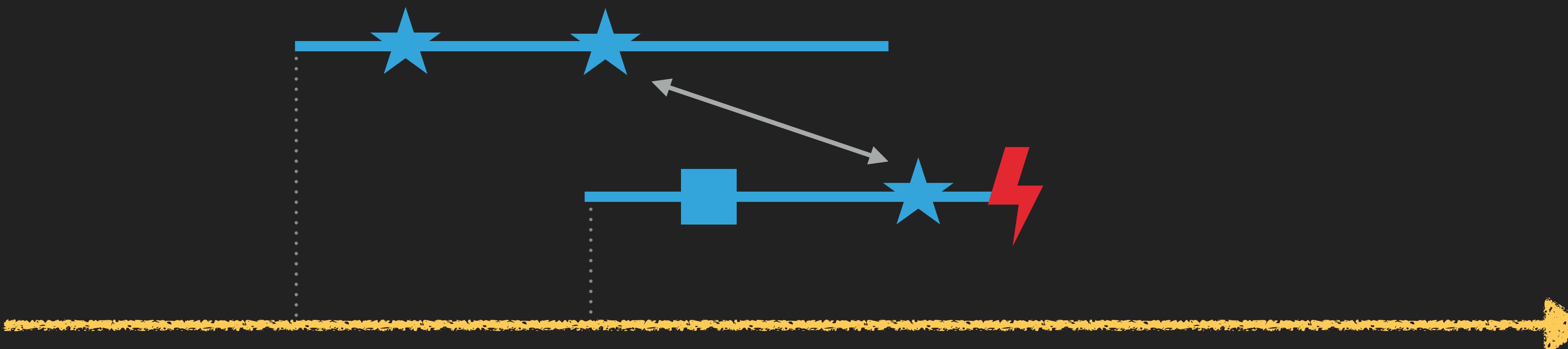
SNAPSHOT ISOLATION



SNAPSHOT ISOLATION



SNAPSHOT ISOLATION



WRITE SKEW

$a, b = 1, 2$

TX1

$t1 = a$

$b = t1$

TX2

$t2 = b$

$a = t2$



$a, b = 2, 1$

WRITE SKEW



TX1

$t_1 = a$
 $b = t_1$

TX2

$t_2 = b$
 $a = t_2$

TX2

$t_2 = b$
 $a = t_2$

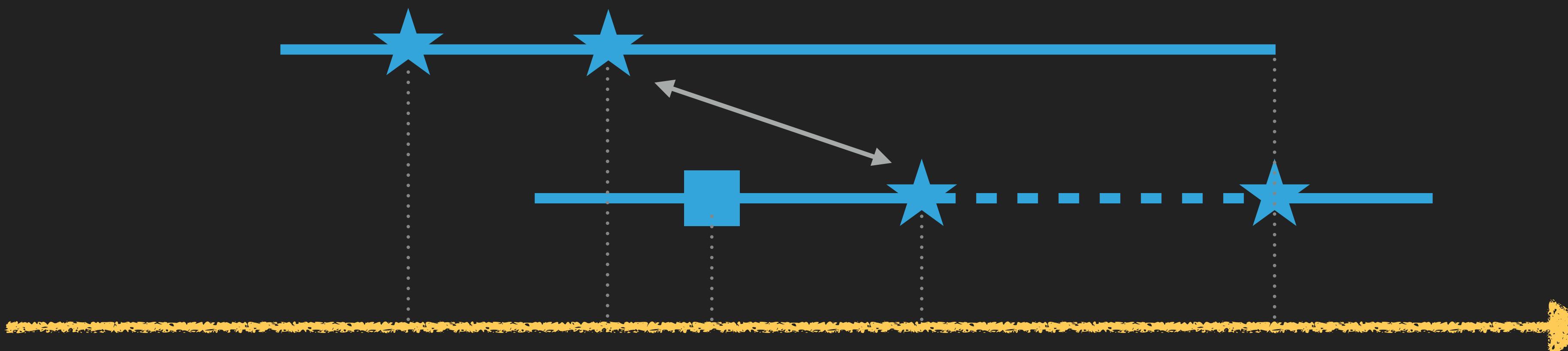
TX1

$t_1 = a$
 $b = t_1$

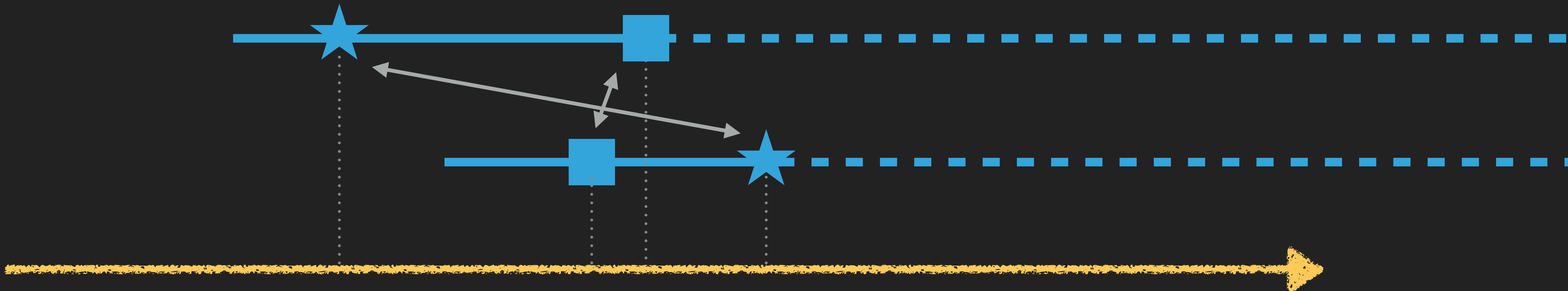
$a, b = 1, 1$

$a, b = 2, 2$

READ COMMITTED



READ COMMITTED

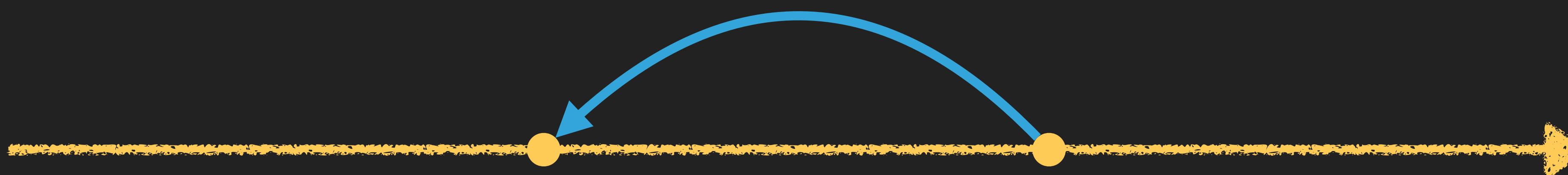


READ COMMITTED

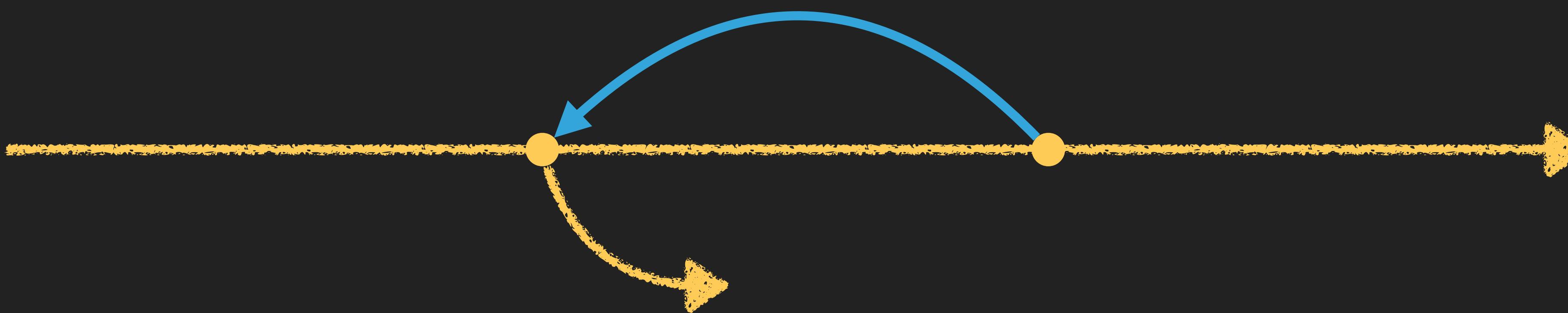
DEADLOCK!



AS OF QUERY



BITEMPORAL MODEL

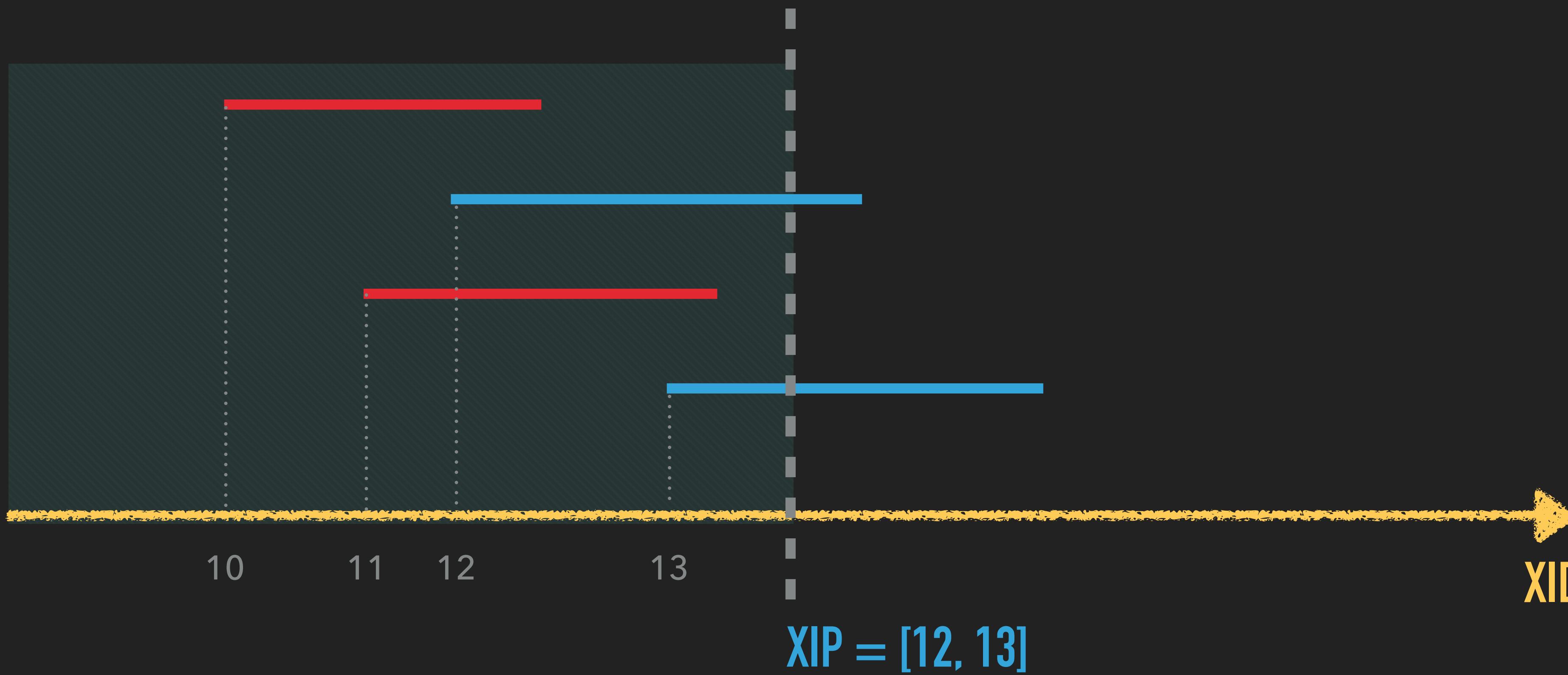


POSTGRES S.I.

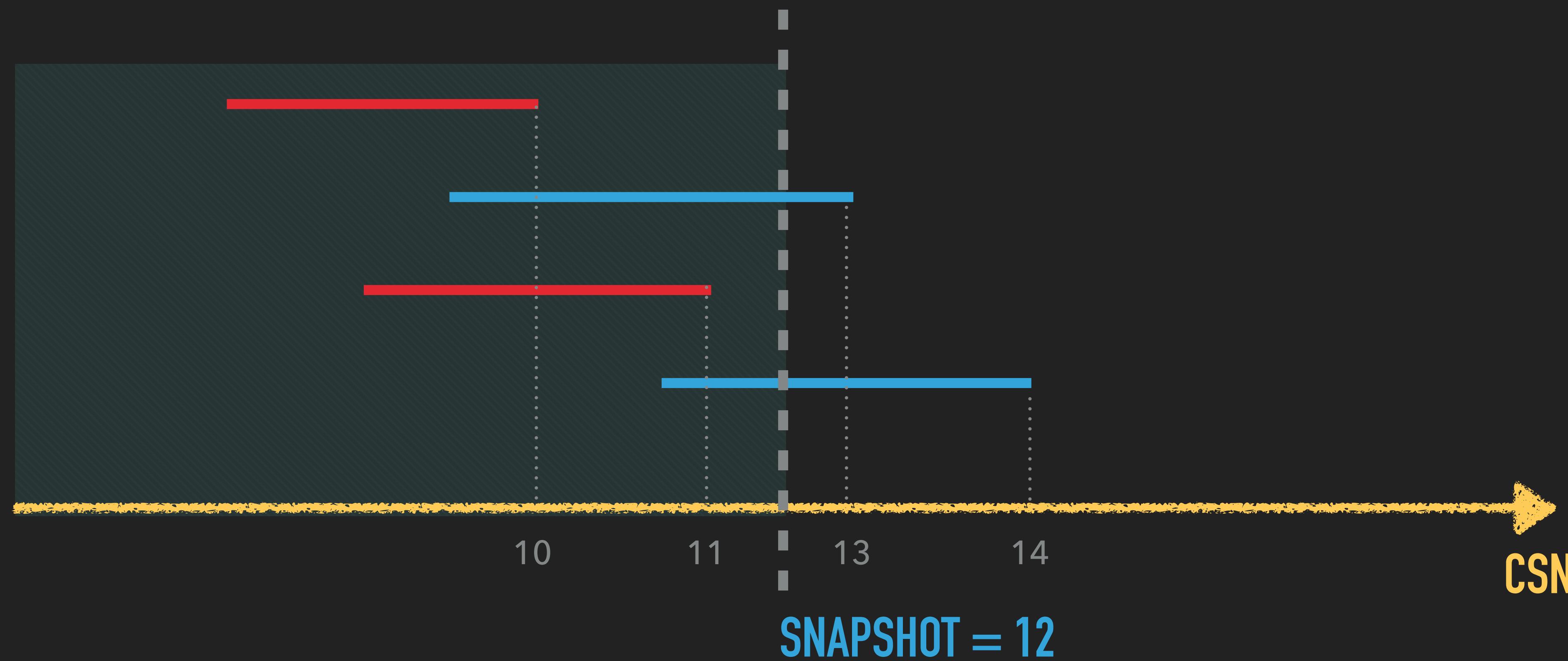
xmin	xmax	key	value
------	------	-----	-------

10	15	foo	41
15	0	foo	42

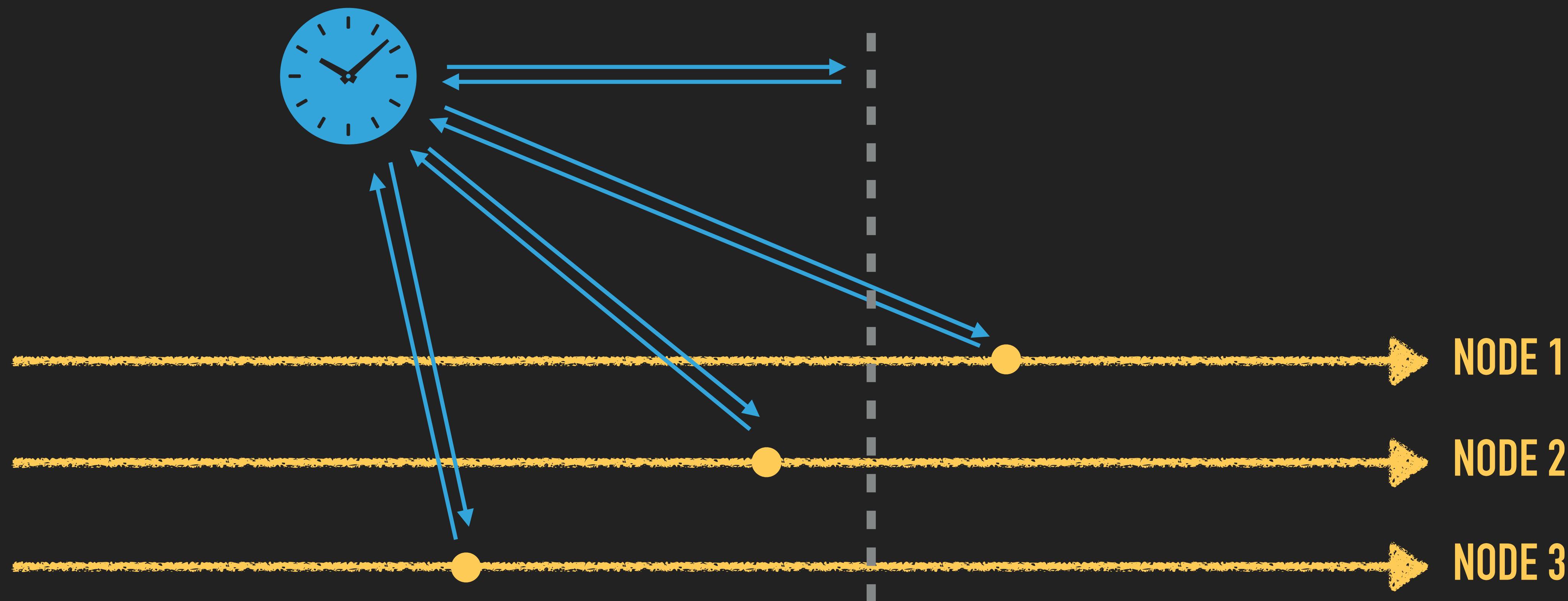
POSTGRES S.I.



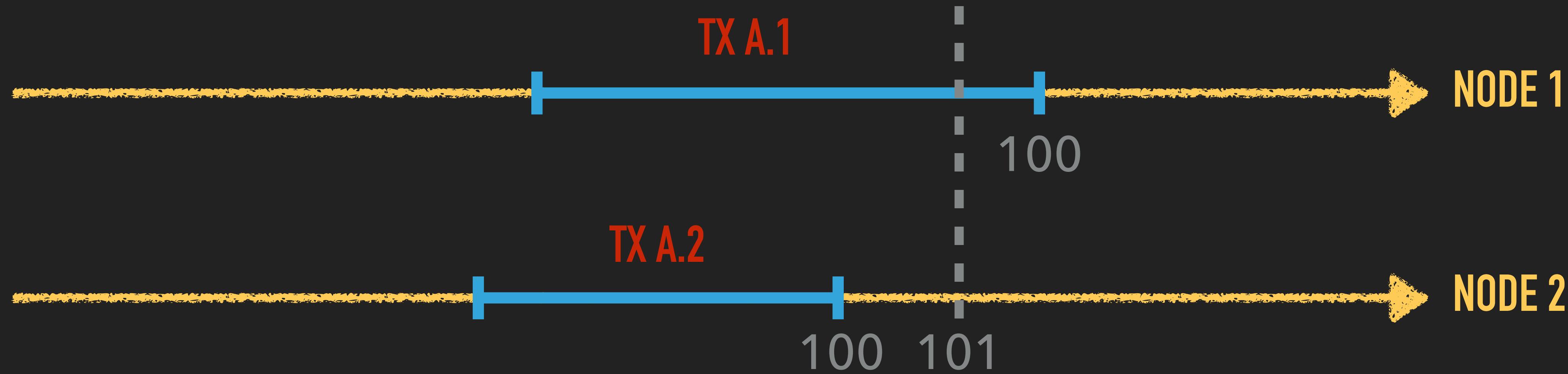
CSN-BASED S.I.



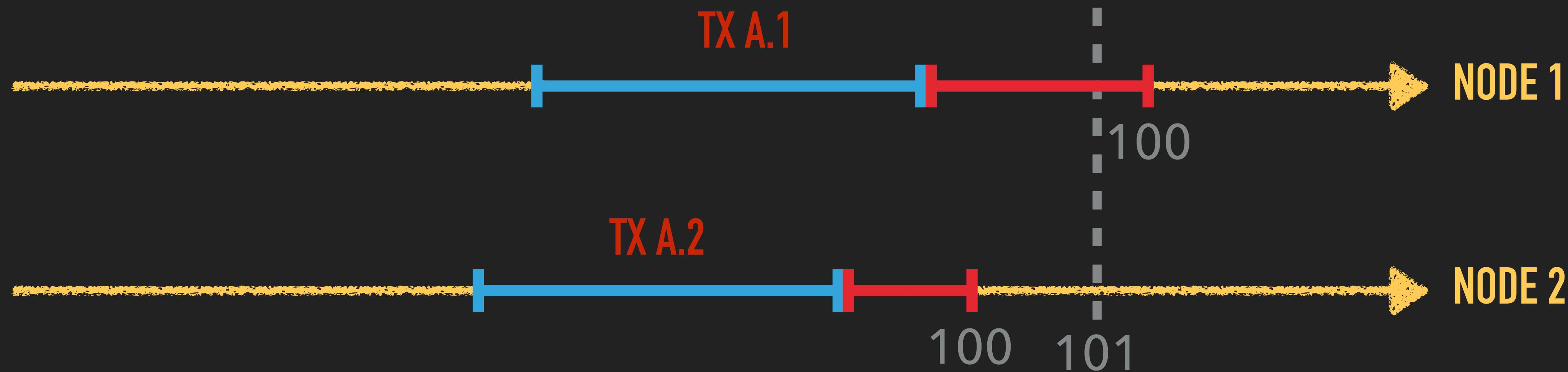
DISTRIBUTED S.I.



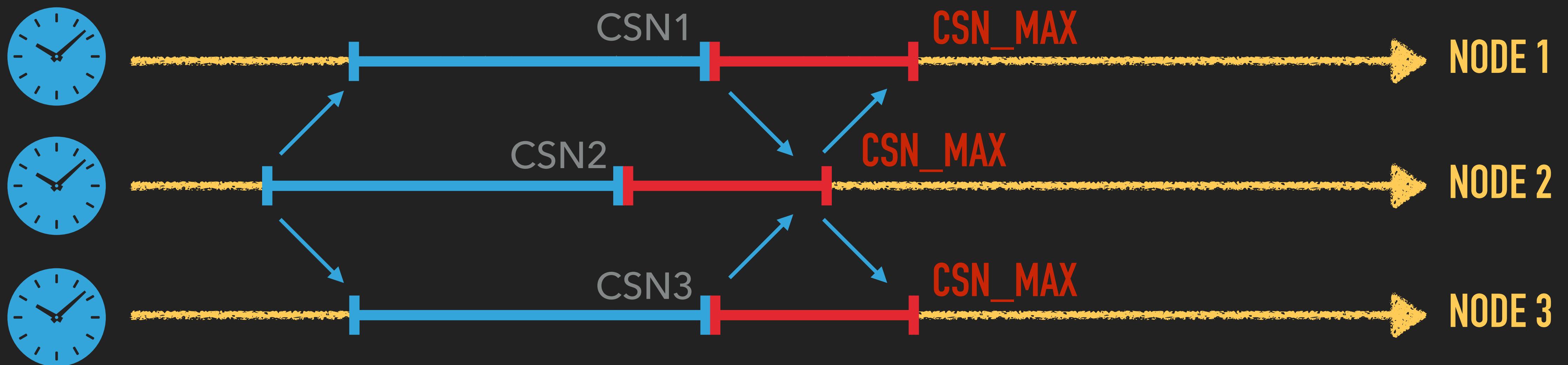
DISTRIBUTED S.I.



DISTRIBUTED S.I.

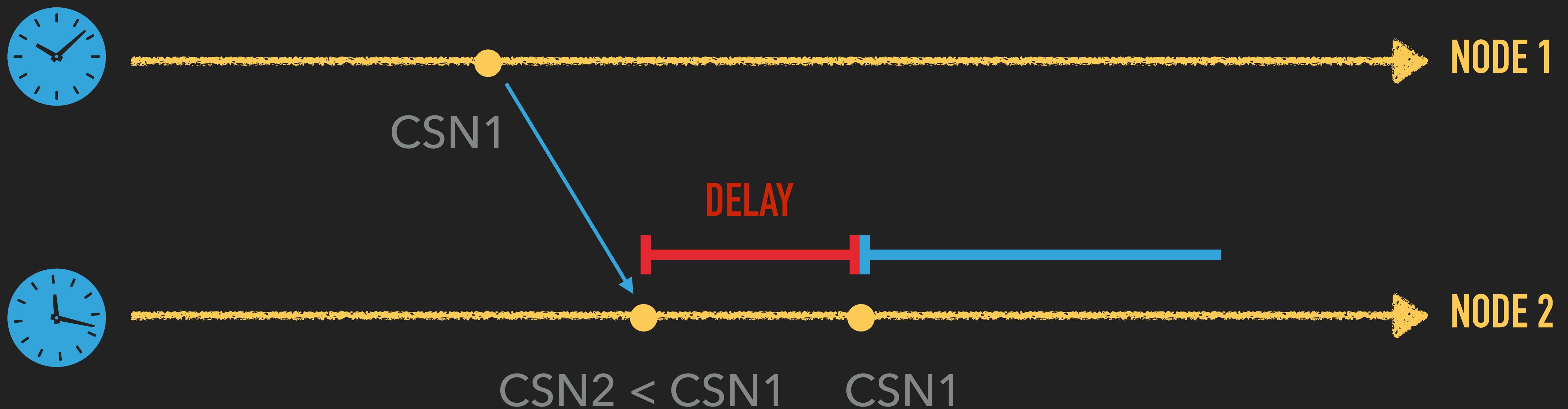


CLOCK-SI

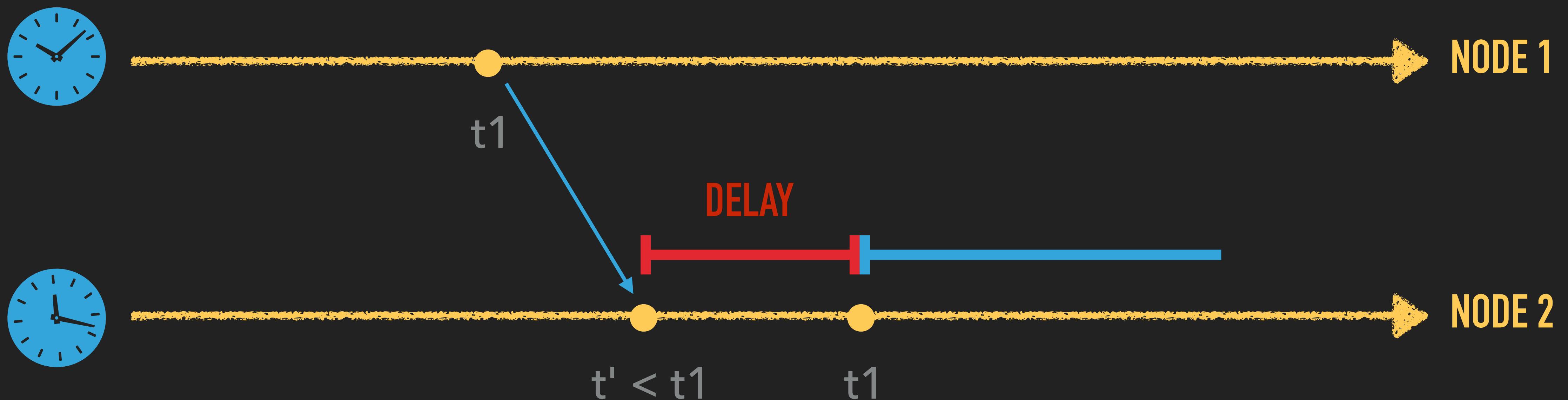


$$\text{CSN_MAX} = \text{MAX}(\text{CSN1}, \text{CSN2}, \text{CSN3})$$

CLOCK-SI: CLOCK SKEW



CLOCK-SI: CLOCK SKEW



ТЕКУЩЕЕ СОСТОЯНИЕ

- ▶ Патч с реализацией Clock-SI для постгреса
- ▶ Патч для postgres_fdw с прозрачной поддержкой Clock-SI
- ▶ Нужно больше верификации
- ▶ Profit

ЧТО ПОЧИТАТЬ

- ▶ Designing data-intensive applications, M. Kleppmann
- ▶ A critique of ANSI SQL isolation levels, H. Berenson
- ▶ Generalized Isolation Level Definitions, A. Adya
- ▶ Clock-SI: Snapshot Isolation for Partitioned Data Stores, Du, Elnikety

ВСЕМ СПАСИБО